



SPLASH

Distributed Systems Made Simple

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Frick, Switzerland

About Me

- 2010-now : PhD student at the University of Neuchatel, Switzerland (and Lua user since then)
 - Topic: large-scale distributed systems, cloud computing. Daily job:
 1. Invent new protocols for cloud applications
 2. Do experiments, write papers, go to I
- 2007-2009: Research engineer at INRIA, France
- BSc and MSc in Computer Engineering at the Università degli Studi Roma Tre, Italy

Motivations

Motivations

- Developing, testing and tuning distributed applications is **hard**
- In Computer Science research, fixing the gap of simplicity between pseudocode description and implementation is **hard**
- Using worldwide testbeds is **hard**



What is PLANETLAB



What is



PLANETLAB



What is



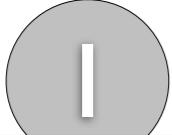
PLANETLAB

- Machines contributed by universities, companies, etc.
 - 1098 nodes at 531 sites (02/09/2011)
 - Shared resources, no privileged access
- University-quality Internet links
- High resource contention
- Faults, churn, packet-loss is the norm
 - Challenging conditions

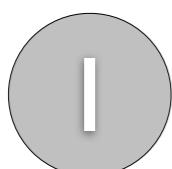


Daily Job With Distributed Systems

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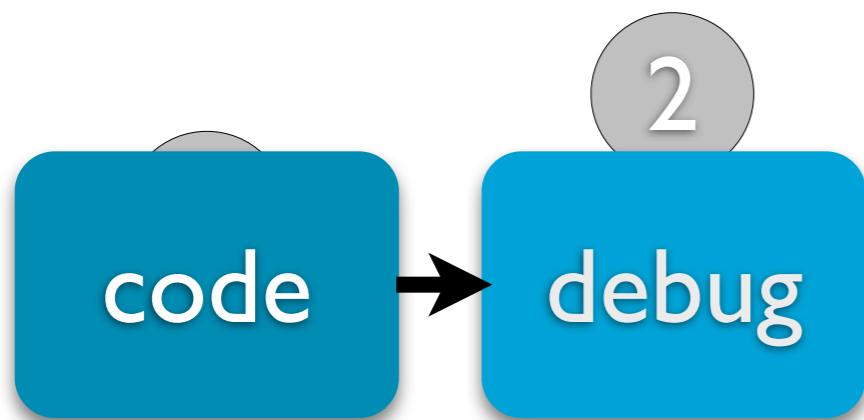


code



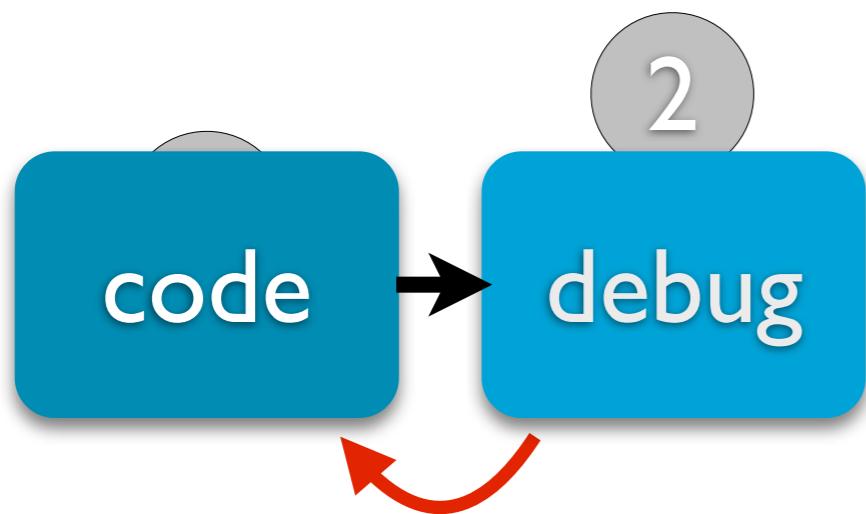
- Write (testbed specific) code
 - Local tests, in-house cluster, PlanetLab...

Daily Job With Distributed Systems



- Debug (in this context, a nightmare)

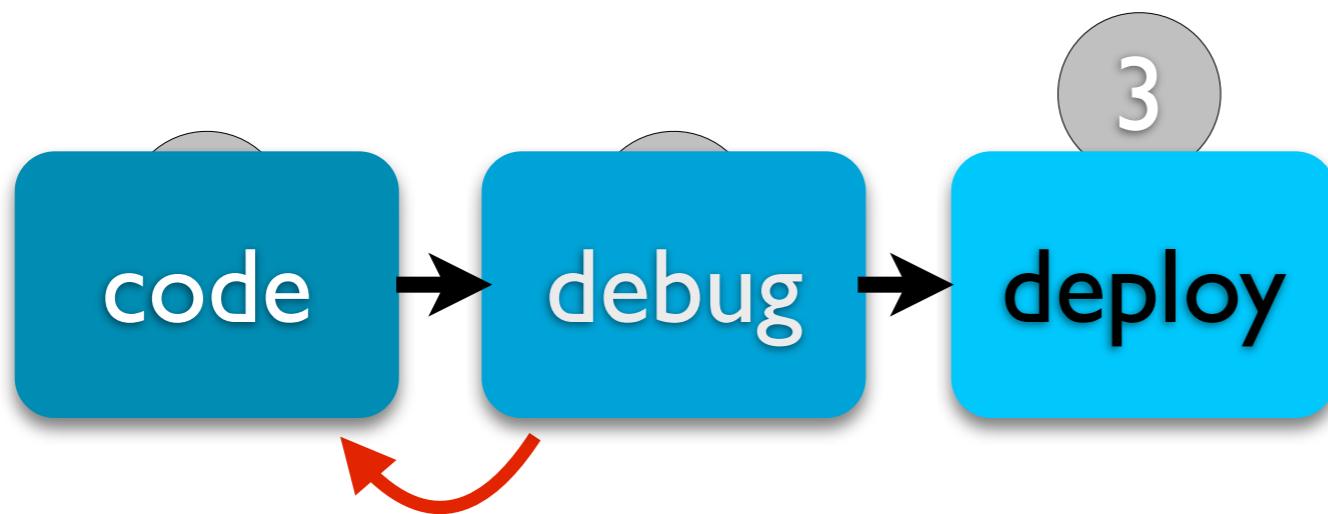
Daily Job With Distributed Systems



2

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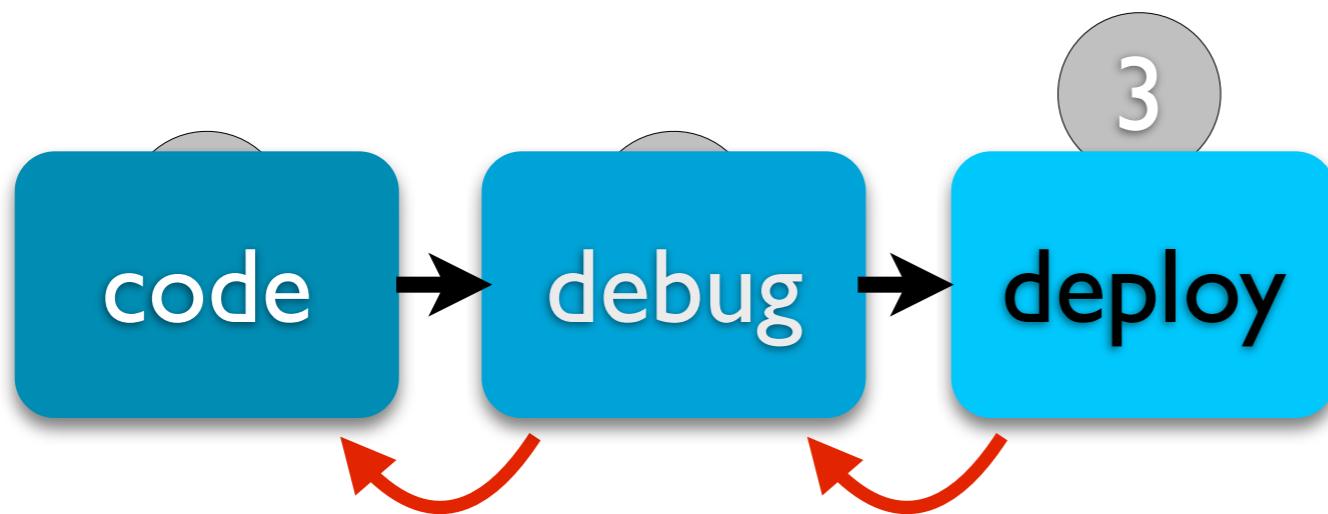
Daily Job With Distributed Systems



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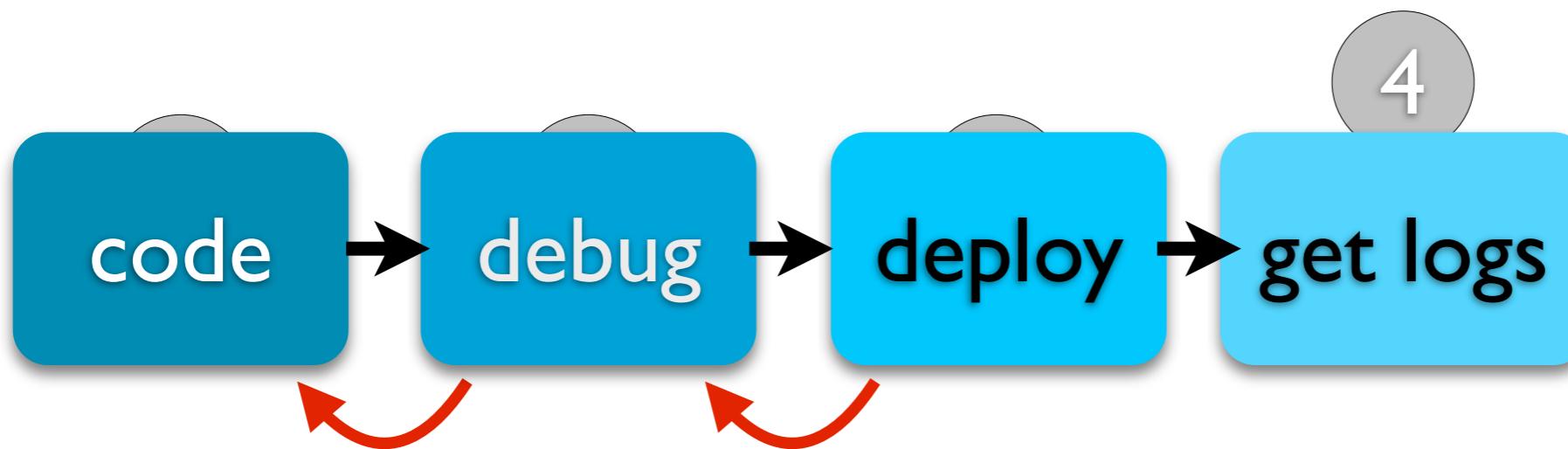
- Deploy, with testbed specific scripts

Daily Job With Distributed Systems



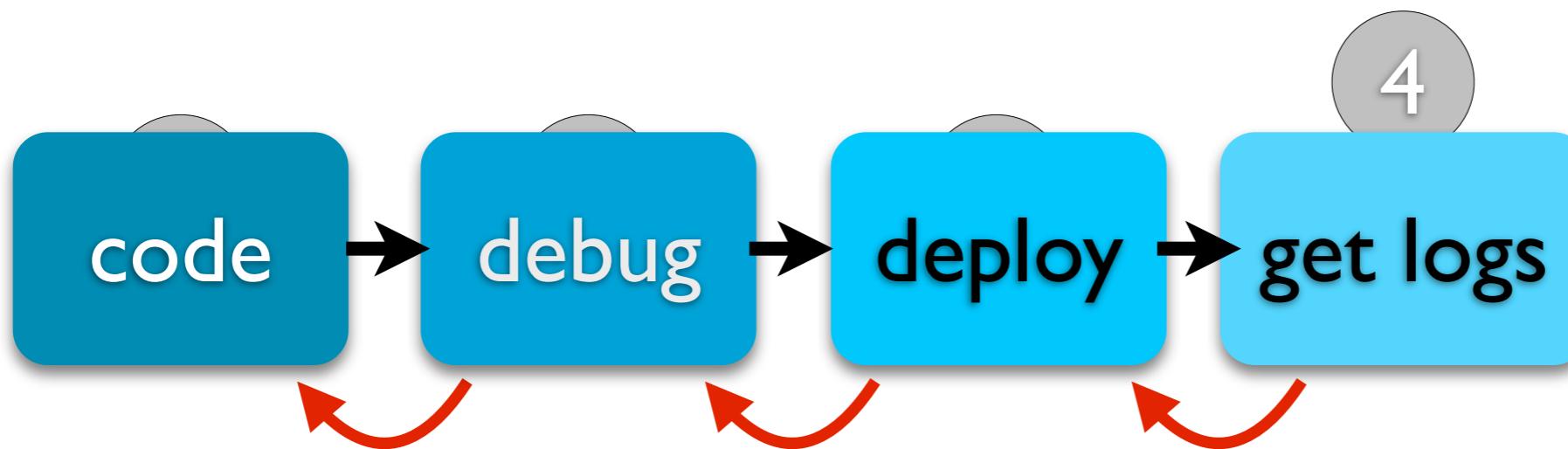
- 3 • Deploy, with testbed specific scripts

Daily Job With Distributed Systems



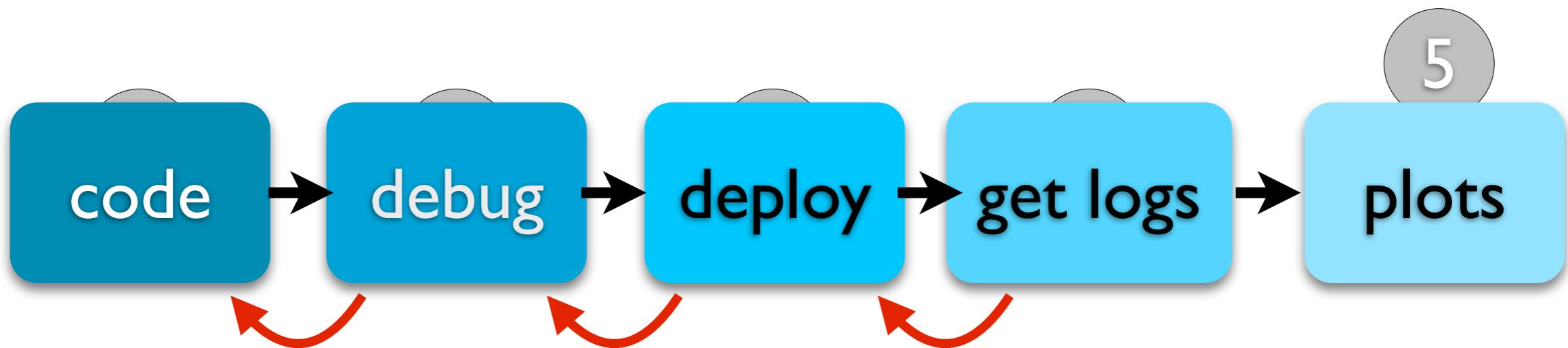
- 4 • Get logs, with testbed specific scripts

Daily Job With Distributed Systems



- 4 • Get logs, with testbed specific scripts

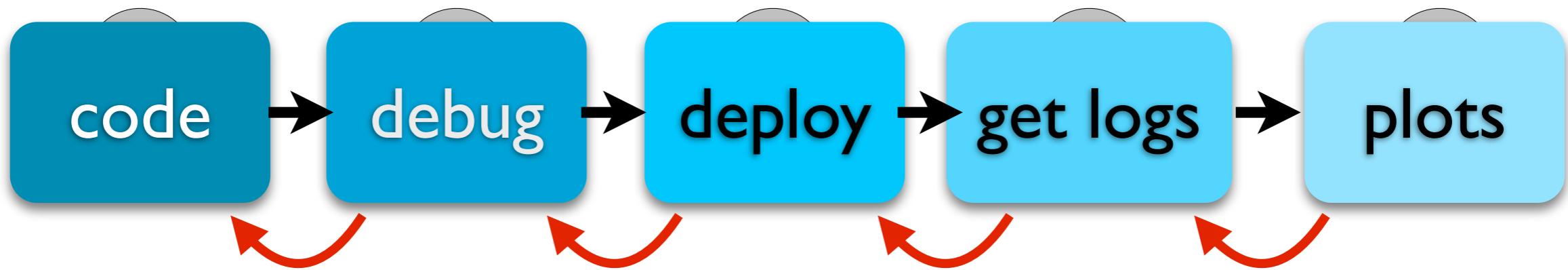
Daily Job With Distributed Systems



5

- Produce plots, hopefully

Daily Job With Distributed Systems



5

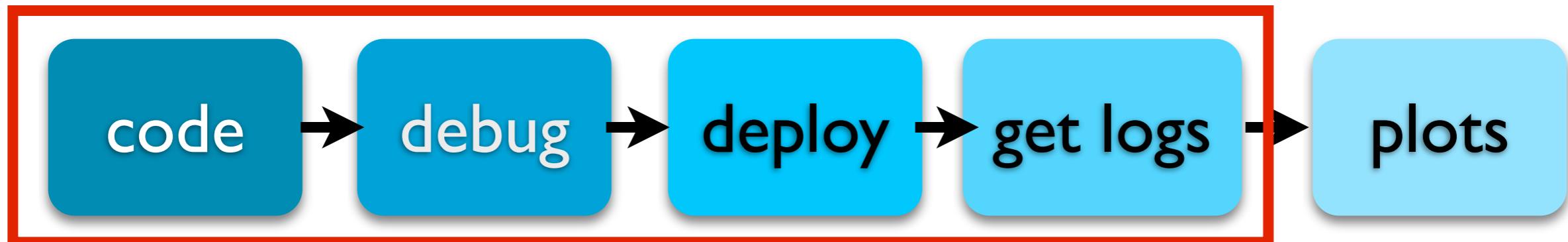
- Produce plots, hopefully

SPLA \downarrow At Glance



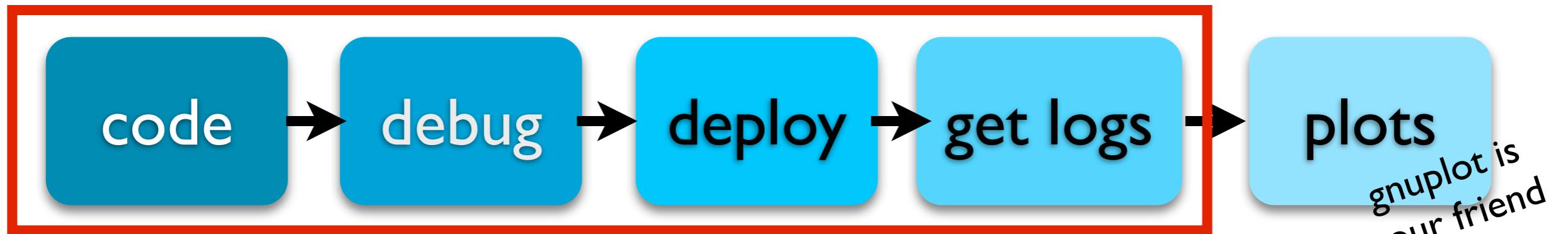
- Supports the **development, evaluation, testing, and tuning** of distributed applications on any testbed:
- In-house cluster, shared testbeds, emulated environments
- Provides an **easy-to-use pseudocode-like language implemented in Lua**

SPLA \downarrow At Glance



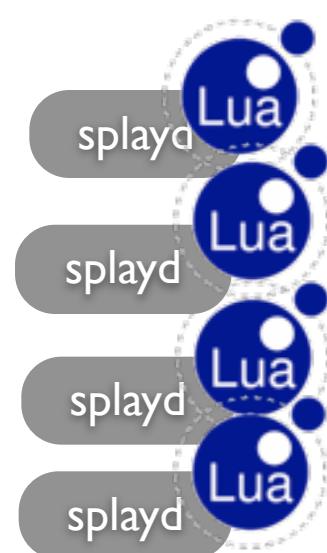
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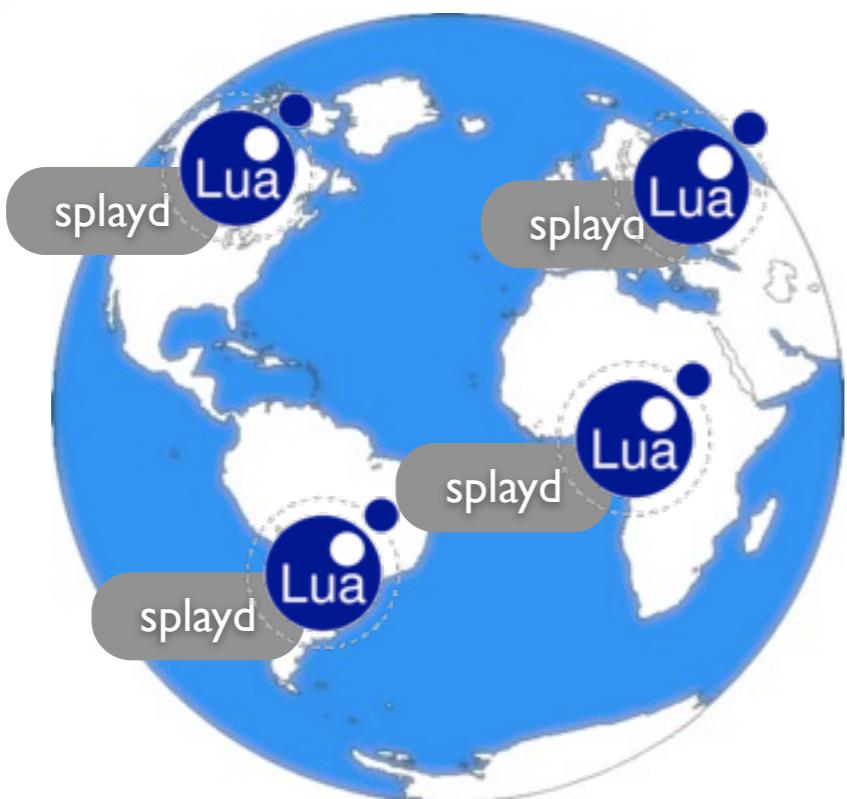
The Big Picture



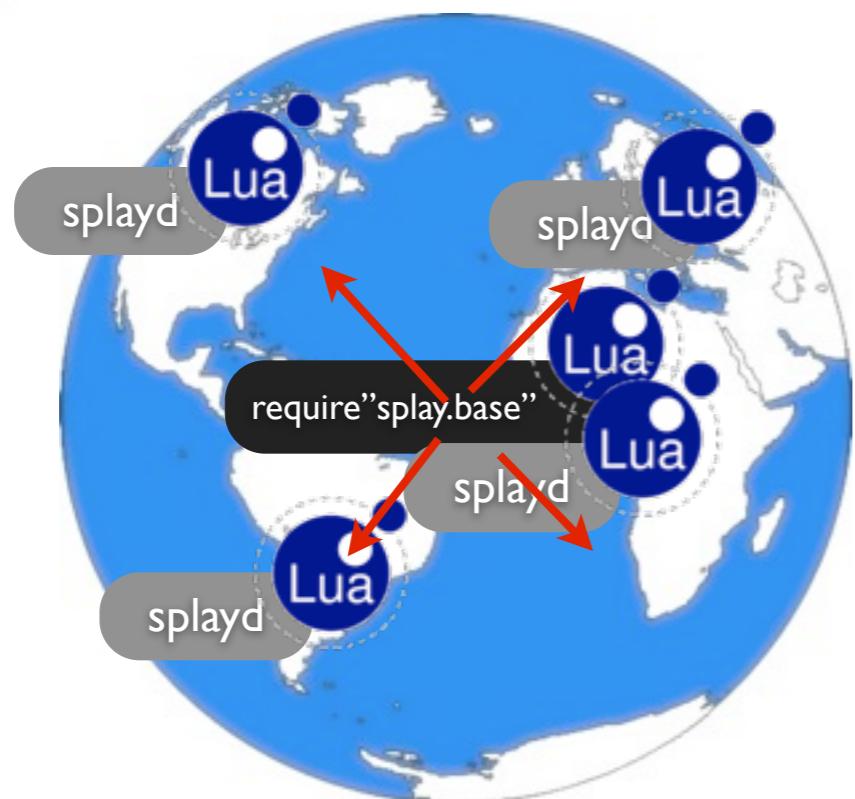
`require "splay.base"`



The Big Picture



The Big Picture



Why



?

Why Lua ?



- Light & Fast
- (Very) Close to equivalent code in C

Why ?

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- Concise
 - Allow developers to focus on ideas more than implementation details
 - Key for researchers

Why ?



- Light & Fast
 - (Very) Close to equivalent code in C
- Concise
 - Allow developers to focus on ideas more than implementation details
 - Key for researchers
- Sandbox thanks to the possibility of easily redefine (even built-in) functions

Concise

```

// ask node n to find the successor of id
n.find_successor(id)
  if (id ∈ (n, successor])
    return successor;
  else
    n' = closest_preceding_node(id);
    return n'.find_successor(id);

// search the local table for the highest predecessor of id
n.closest_preceding_node(id)
  for i = m downto 1
    if (finger[i] ∈ (n, id))
      return finger[i];
  return n;

// create a new Chord ring.
n.create()
  predecessor = nil;
  successor = n;
  ...

// join a Chord ring containing node n'.
n.join(n')
  predecessor = nil;
  successor = n'.find_successor(n);

// called periodically, verifies n's immediate
// successor, and tells the successor about n.
n.stabilize()
  x = successor.predecessor;
  if (x ∈ (n, successor))
    successor = x;
  successor.notify(n);

// n' thinks it might be our predecessor.
n.notify(n')
  if (predecessor is nil or n' ∈ (predecessor, n))
    predecessor = n';

// called periodically, refreshes finger table entries.
// next stores the index of the next finger to fix.
n.fix_fingers()
  next = next + 1;
  if (next > m)
    next = 1;
  finger[next] = find_successor(n + 2^{next-1});

// called periodically, checks whether predecessor has failed.
n.check_predecessor()
  if (predecessor has failed)
    predecessor = nil;

```

Concise

Pseudo code
as published
on original
paper

Executable
code using →
SPLAY
libraries

```

require "splay.base"
rpc = require "splay.rpc"
between, call, thread, ping = misc.between_c, rpc.call, events.thread, rpc.ping
n, predecessor, finger, timeout, m = {}, nil, {}, 5, 24
function join(n0) -- n0: some node in the ring
  predecessor = nil
  finger[1] = call(n0, {'find_successor', n0.id})
  call(finger[1], {'notify', n0})
end
function closest_preceding_node(id)
  for i = m, 1, -1 do
    if finger[i] and between(finger[i].id, n.id, id) then
      return finger[i]
    end
  end
  return n
end
function find_successor(id)
  if finger[1].id == n.id or between(id, n.id, (finger[1].id + 1) % 2^m) then
    return finger[1]
  else
    local n0 = closest_preceding_node(id)
    return call(n0, {'find_successor', id})
  end
end
function stabilize()
  local x = call(finger[1], 'predecessor')
  if x and between(x.id, n.id, finger[1].id) then
    finger[1] = x -- new successor
    call(finger[1], {'notify', n0})
  end
end
function notify(n0)
  if n0.id == n.id and
     (not predecessor or between(n0.id, predecessor.id, n.id)) then
    predecessor = n0
  end
end
function fix_fingers()
  refresh = (refresh and (refresh % m) + 1) or 1 -- 1 ≤ next ≤ m
  finger[refresh] = find_successor((n.id + 2^(refresh - 1)) % 2^m)
end
function check_predecessor()
  if predecessor and not rpc.ping(predecessor) then
    predecessor = nil
  end
end
n.id = math.random(1, 2^m)
finger[1] = n
if job then
  n.ip, n.port = job.me.ip, job.me.port
  thread(function() join({ip = "192.42.43.42", port = 20000}) end)
else
  n.ip, n.port = "127.0.0.1", 20000
  if arg[1] then n.ip = arg[1] end
  if arg[2] then n.port = tonumber(arg[2]) end
  if not arg[3] then
    print("RDV")
  else
    print("JOIN")
    thread(function() join({ip = arg[3], port = tonumber(arg[4])}) end)
  end
end

```

Sandbox: Motivations

- Experiments should access only their own resources
- Required for non-dedicated testbeds
 - In universities, companies
 - Totally available at night/holiday time
- Memory-allocation, filesystem, network resources are restricted



Sandboxing LuaSocket

- Same API as plain LuaSocket
- On-demand sandboxed sockets
 - Very easy thanks to Lua's metatable
- Limits chosen by the SPLAYd deployer
- Both TCP and UDP
- Example: UDP's socket.send

Sandboxing LuaSocket



```
if sock.send then
    new_sock.send = function(self, data)
        local len = #data
        if total_sent + len > max_send then
            l_o:warn("Send restricted (total: "..total_sent..")")
            return nil, "restricted"
        end
        local n, status
        if math.random(1000) > udp_drop_ratio then
            n, status = sock:send(data)
        else
            n = len
        end
        if n then
            total_sent = total_sent + len
        end
        return n, status
    end
end
```

```
for _,module in pairs(splay)
```

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```

luasocket

events

io

crypto

llenc/benc

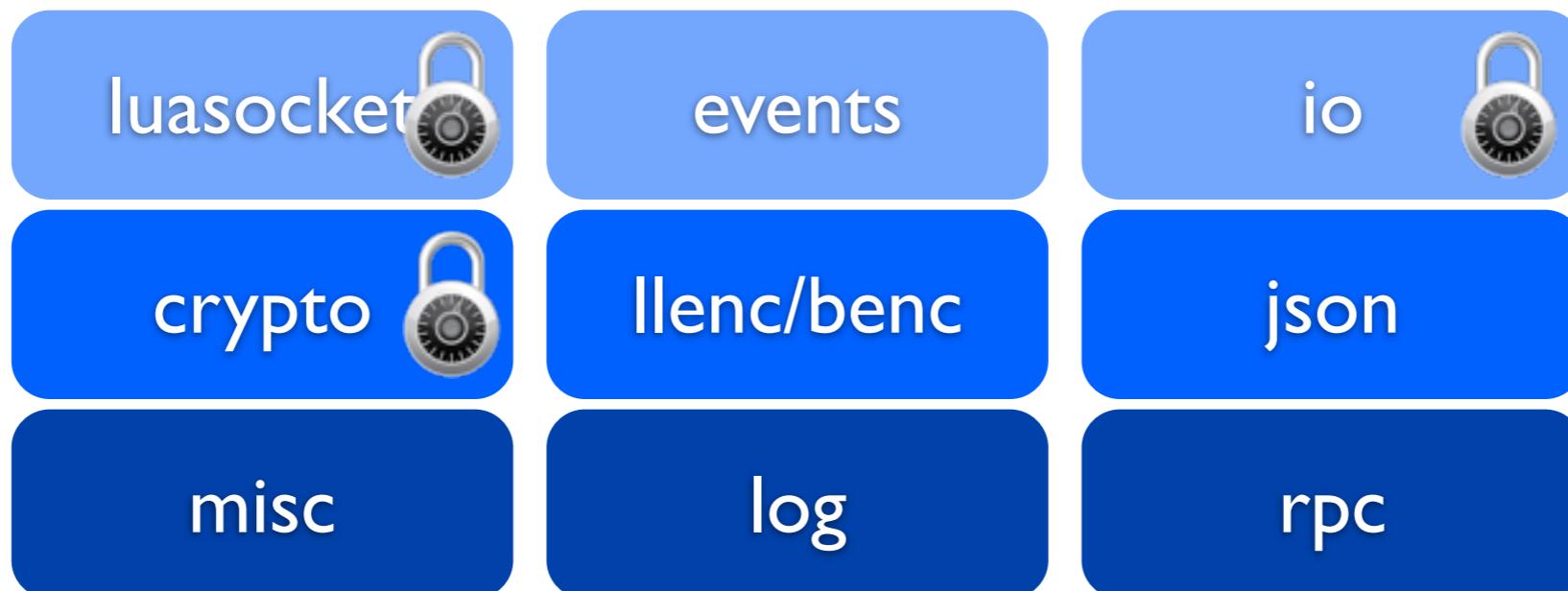
json

misc

log

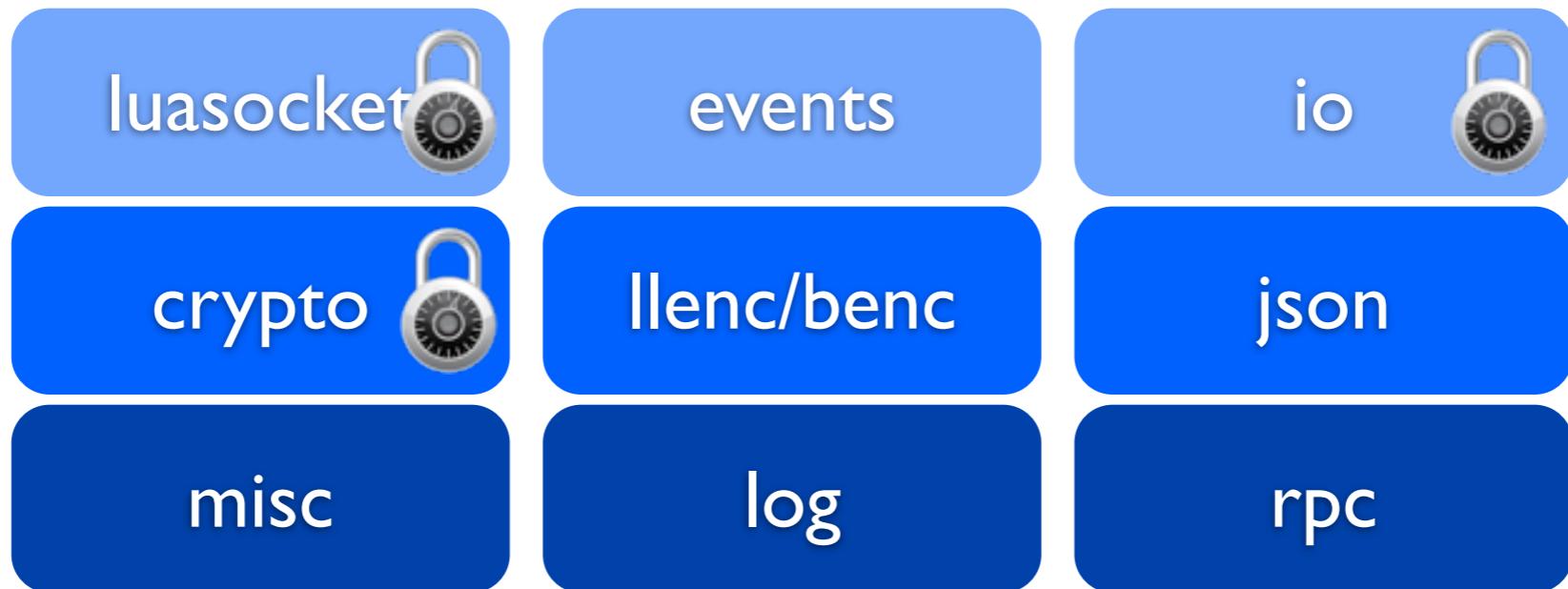
rpc

```
for _,module in pairs(splay)
```



Modules sandboxed to prevent
access to sensible resources

splay.events



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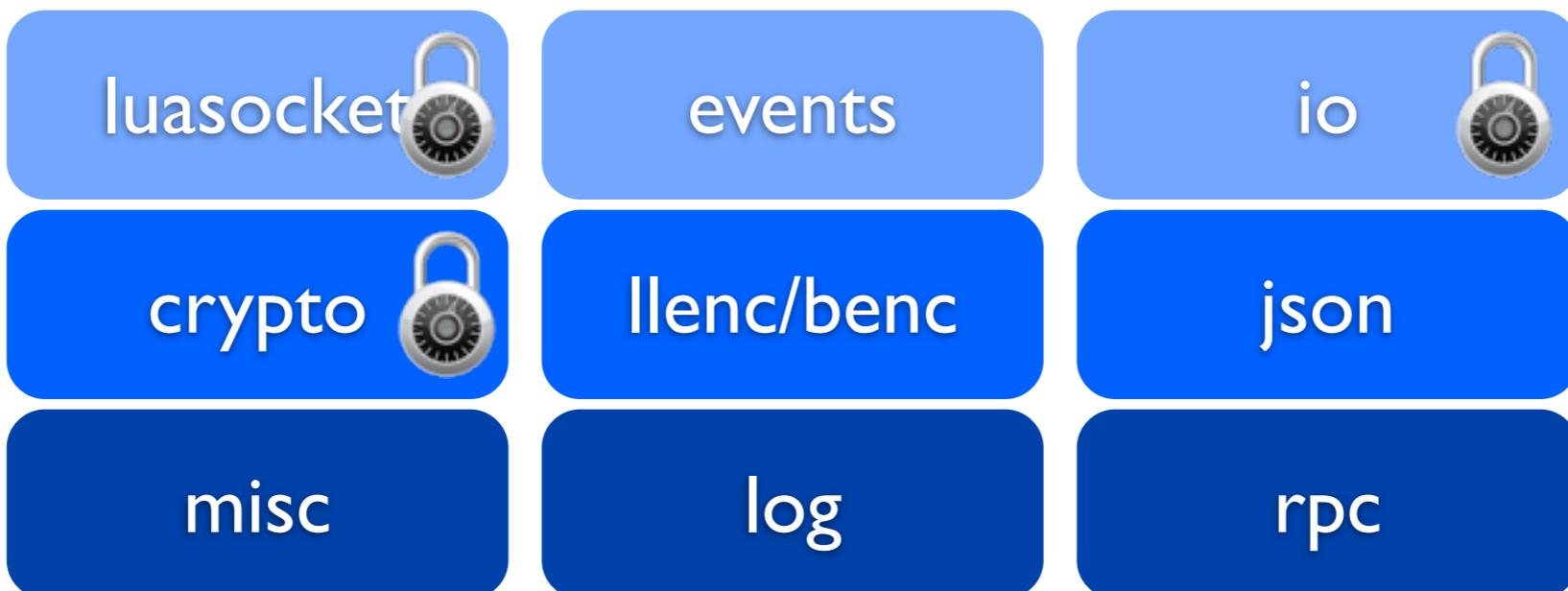
events

splay.events

- Distributed protocols can use message-passing paradigm to communicate
- Nodes react to events
 - Local, incoming, outgoing messages
- The **core** of the Splay runtime
 - `splay.socket`, `splay.io` designed to provide a non-blocking operational mode
- Based on Lua's coroutines

events

splay.rpc



splay.rpc

rpc

splay.rpc

- Default mechanism to communicate between nodes
- Support for UDP/TCP
- Efficient BitTorrent-like encoding
- Experimental binary encoding

rpc

Life Before SPLAʌ

- Time spent on developing testbed specific protocols
- Or fallback on simulations...
- The focus should be on **ideas**
- Researchers usually have no time to produce industrial-quality code

```

public void declareNeighbourInactive(final DHTNode node) {
    // DO NOTHING HERE:
}

public void addToLeafSet(final DHTNode node) {
    // DO NOTHING
}

/*
 * Empty UDP message is 64 bytes, to which we must add 1 bytes for the header,
 * 1 bytes for the messagId, 28 bytes (IP,port,Bamboo ID) for message src, 28
 * bytes for message dest + bytes for the specific data carried in the
 * message. NOTE: the size of informations to ack a message is 10 bytes,
 */
@Override
public long calculateMessageBytes(final Message msg) {
    long result = 122; // 64 + 1 + 1 + 28 + 28;

    if (msg.header == MessageType.PING) {
        /* nothing to add */
        return result;
    }

    if (msg.header == MessageType.PONG) {
        /* some bytes for the informations about which message is being acked */

        result += ACK_INFORMATION_SIZE;
    }

    return result;
}

if (msg.header == MessageType.BAMBOO_JOIN_LOOKUP_REQUEST

```

Life Before SPLASH

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```

    /*

    * CAN BE COMMENTED AS SEND RETRIAL OF MESSAGES IS OFF else { it is a

    * duplicate lookup response, send a PONG to (hopefully) stop source to

    * send us again that message final Message pong = new

    * Message(MessageType.PONG, msg.messageId, this, msg.sender, msg.ackedMsg

    * = msg; simulator.send(pong); }

    */

}

/* ELSE IF IS SOMETHING NOT A LOOKUP RESPONSE */

else {

final Message beingAcked = msg.ackedMsg;

final ResponseArrivedCallback handler = this.sentMessagesCallbacks

.remove(beingAcked);

/* if there is a mapping */

if (handler != null) {

    handler.onResponseReceived(msg);

}

/* else { this may happen in case of a response arrive later than

 * expected... log.info("%%% " + this + " received the response => " +

 * msg +

 * " but no handler was expecting it! So, simply send back a PONG to make it stop sending again..\n"

 * ); send a PONG to (hopefully) stop source to send us again final

 * Message pong = new Message(MessageType.PONG, msg.messageId, this,

msg.sender, msg.ackedMsg = msg; simulator.send(pong); }

*/
}

```

there is more :-)

```

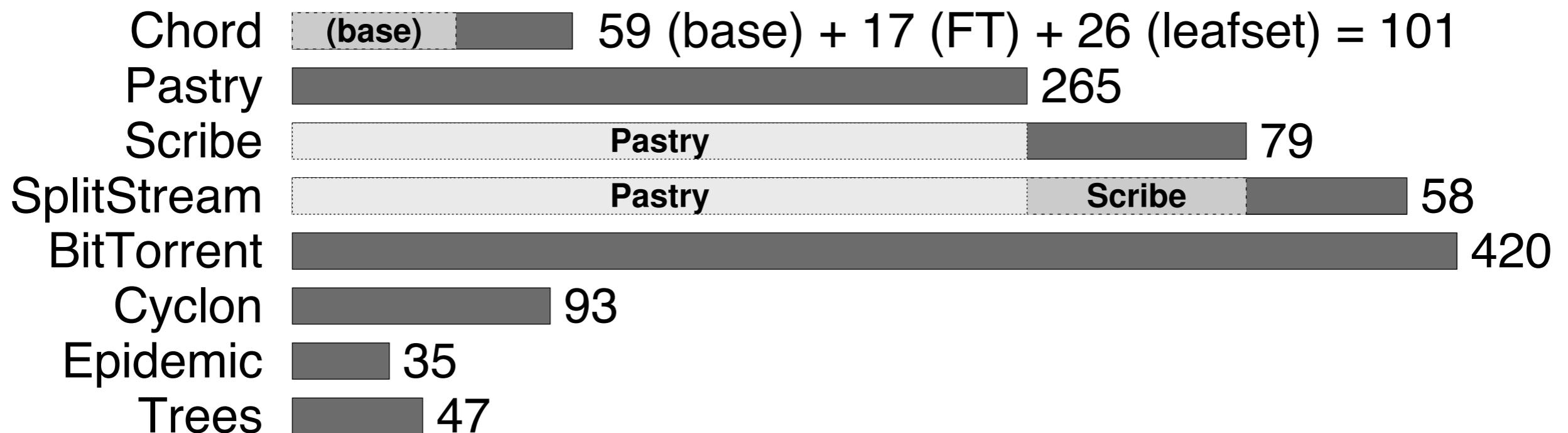
}

// /* it replied at last, so "reset" the value */

```

- Time spent on developing testbed specific protocols
- Or fallback on simulations...
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Life With SPLAʌ



- Lines of pseudocode \approx Lines of executable code

Live Demo



www.splay-project.org



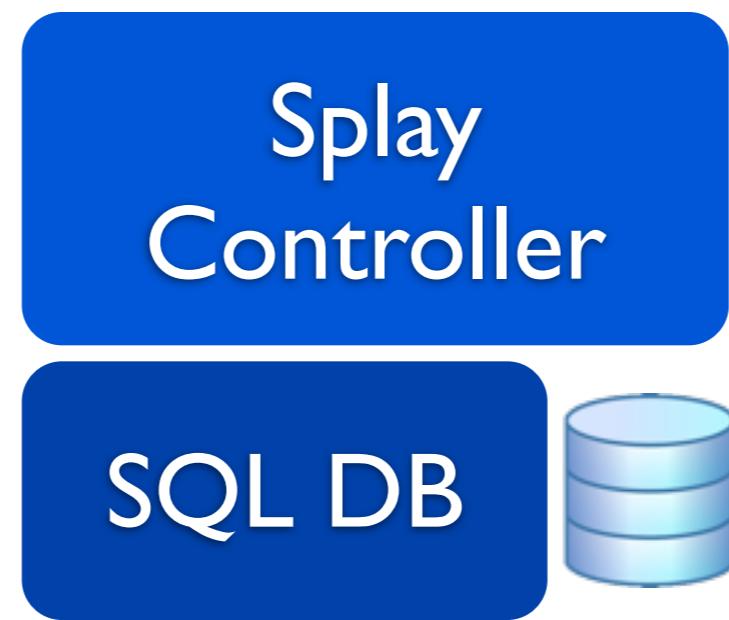
Take-away Slide

- Distributed systems raise a number of issues related to their evaluation
- Their implementation, debug, deployment and tuning is hard
- SPLΛY leverages Lua to produce an easy to use yet powerful working environment to solve these issues

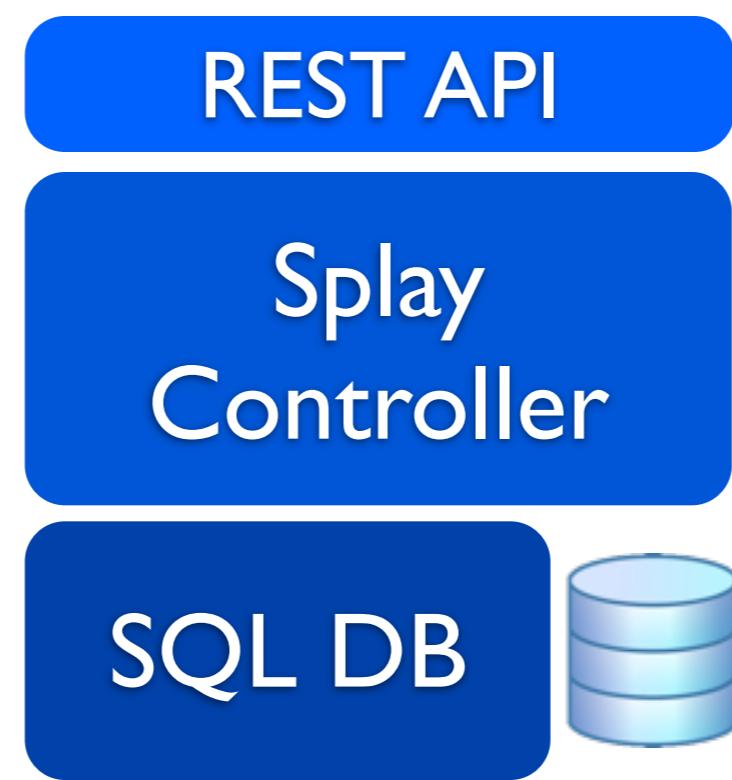
Backup Slides

The Big Picture

The Big Picture



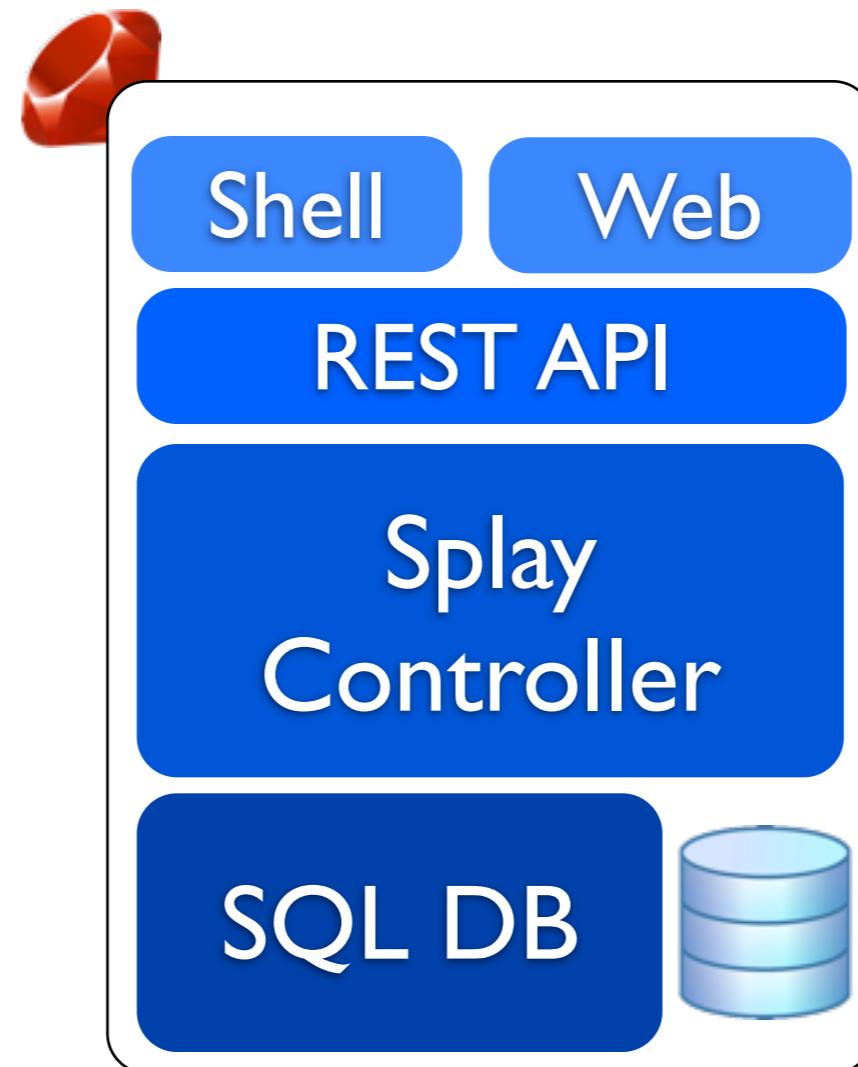
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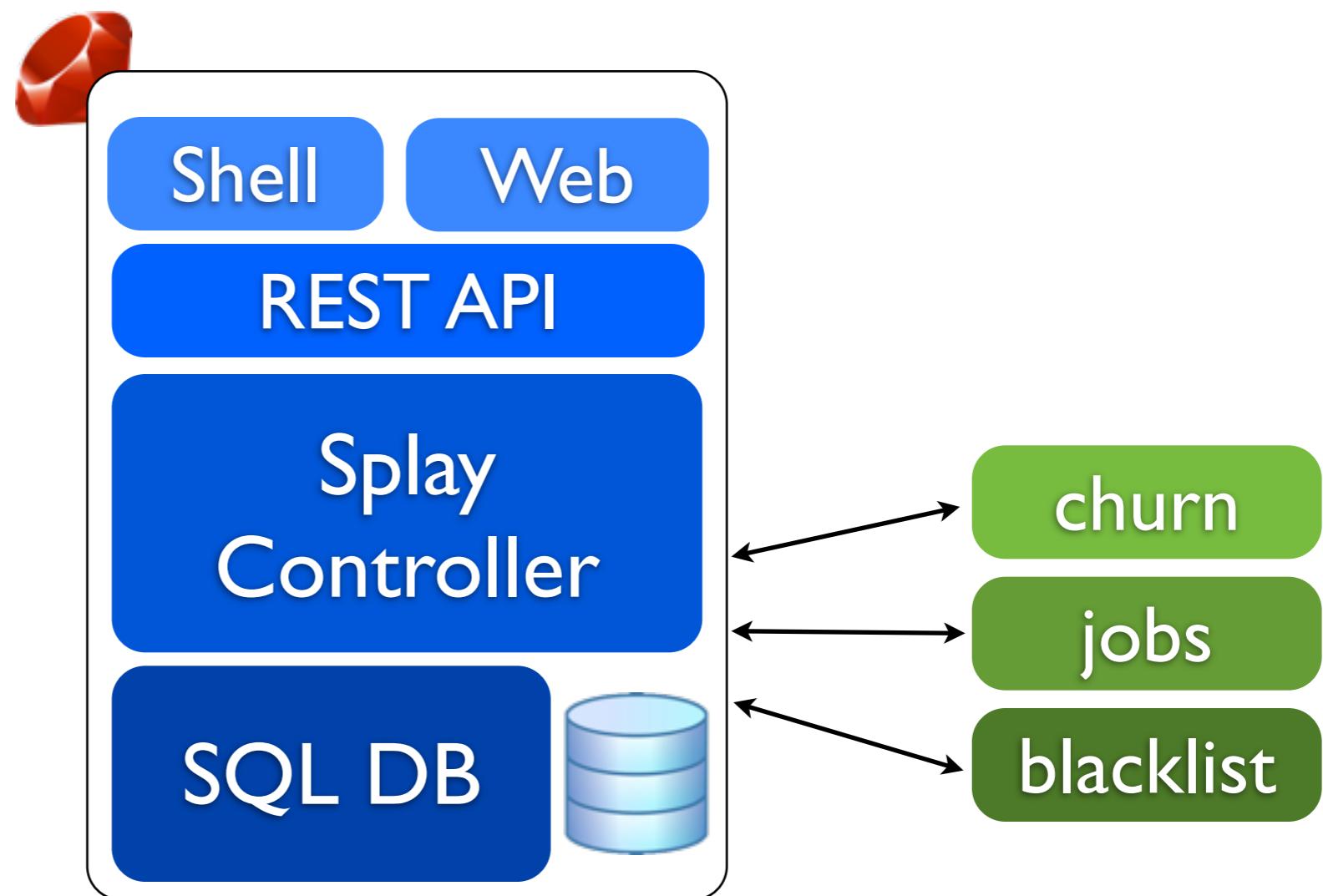
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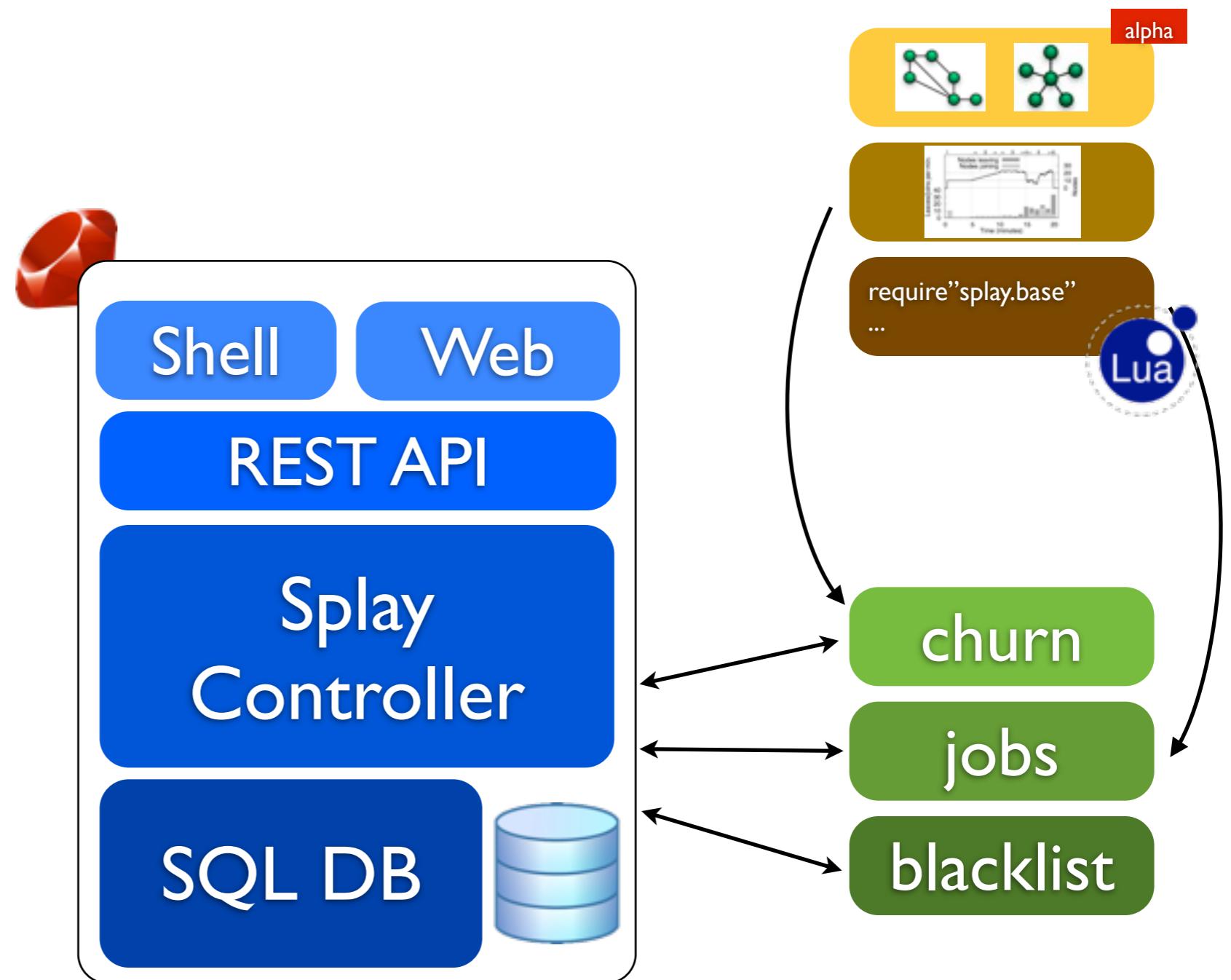
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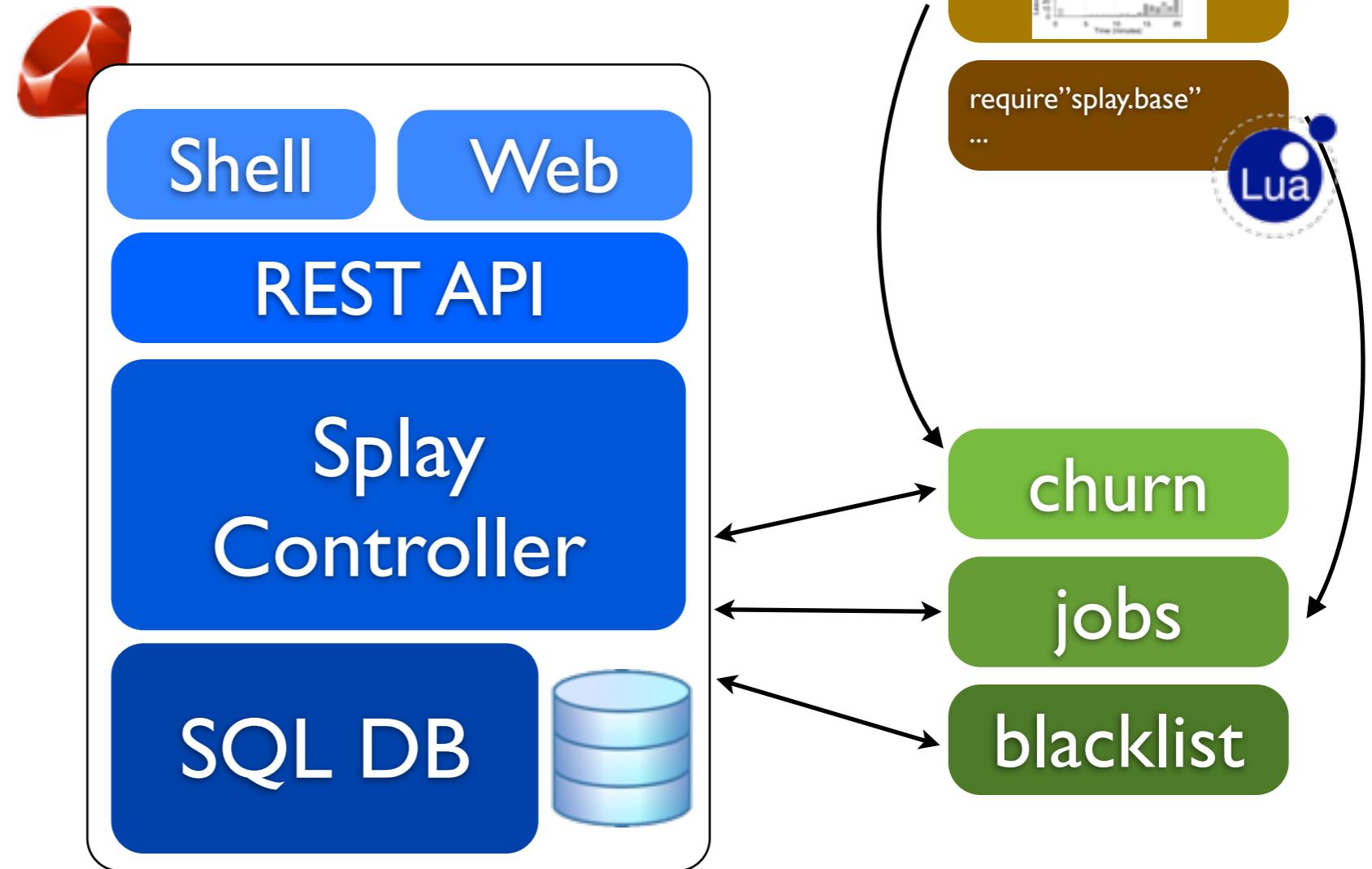
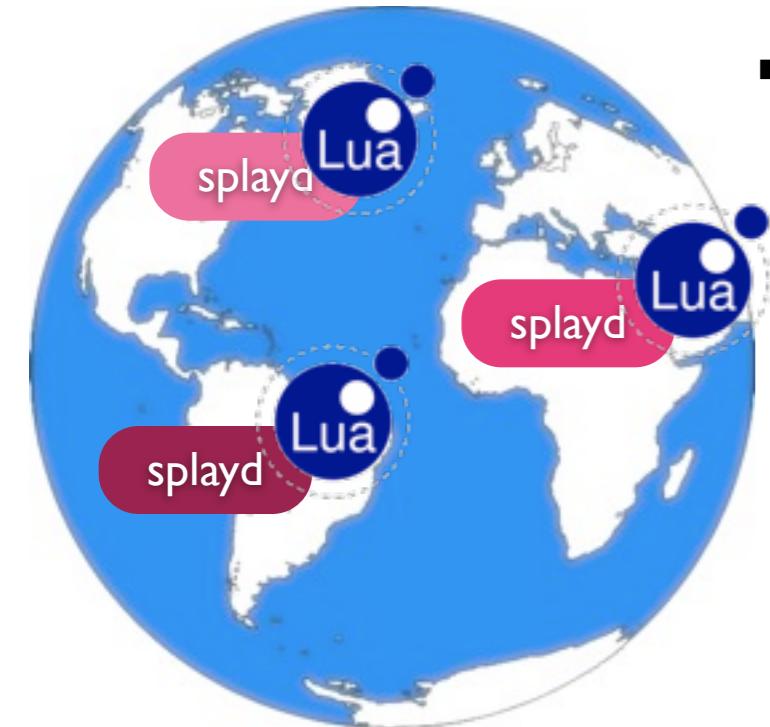
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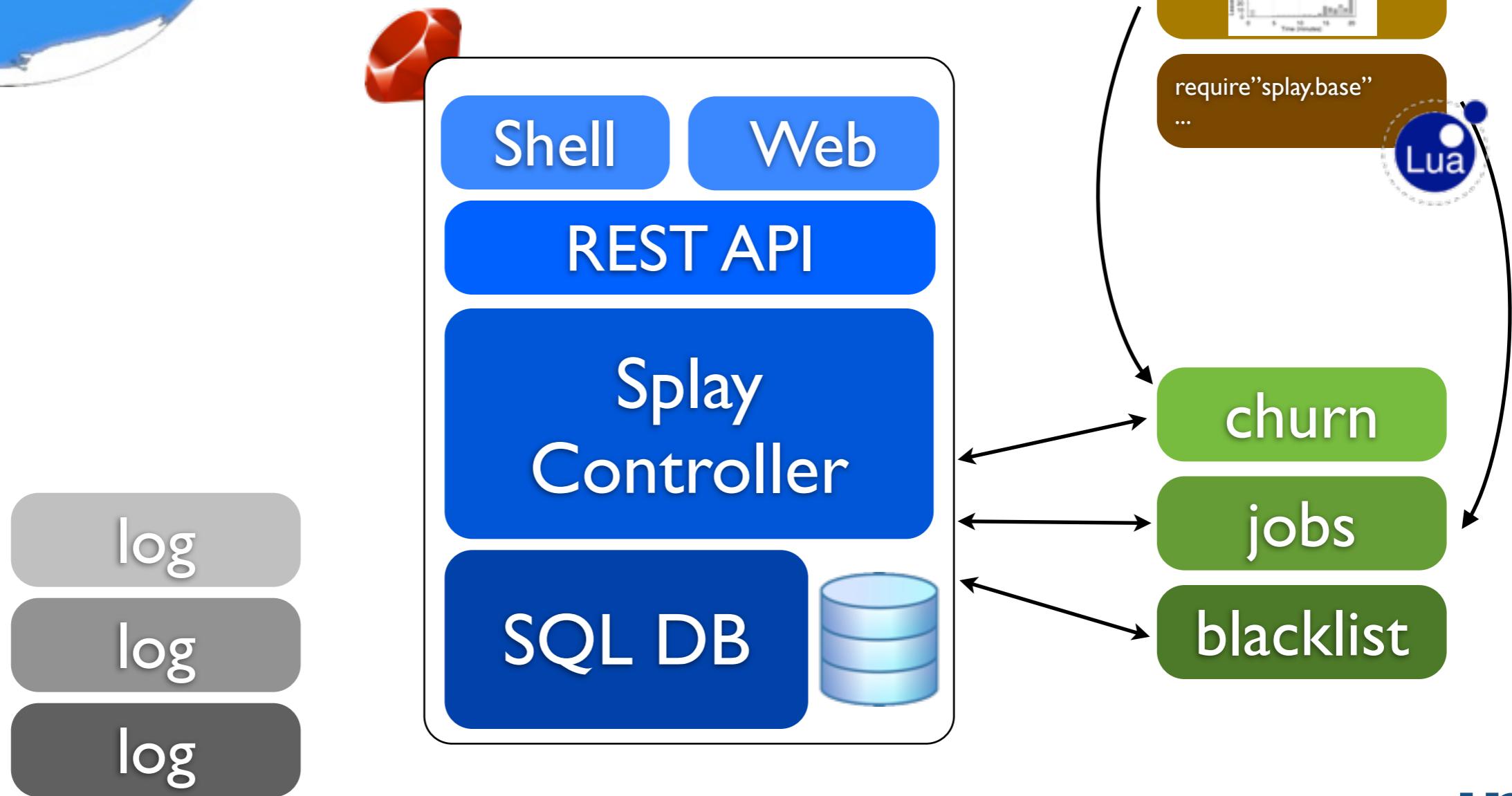
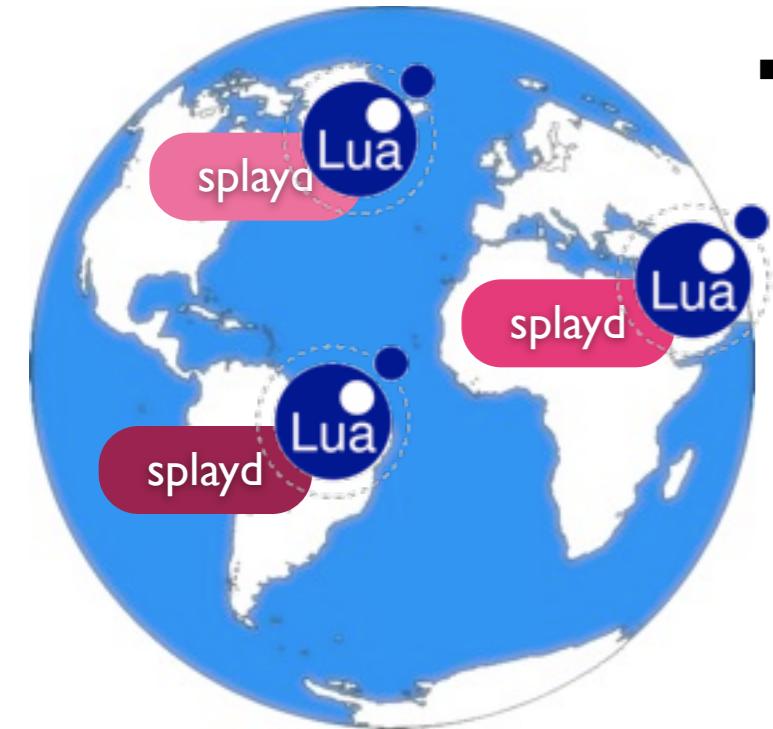
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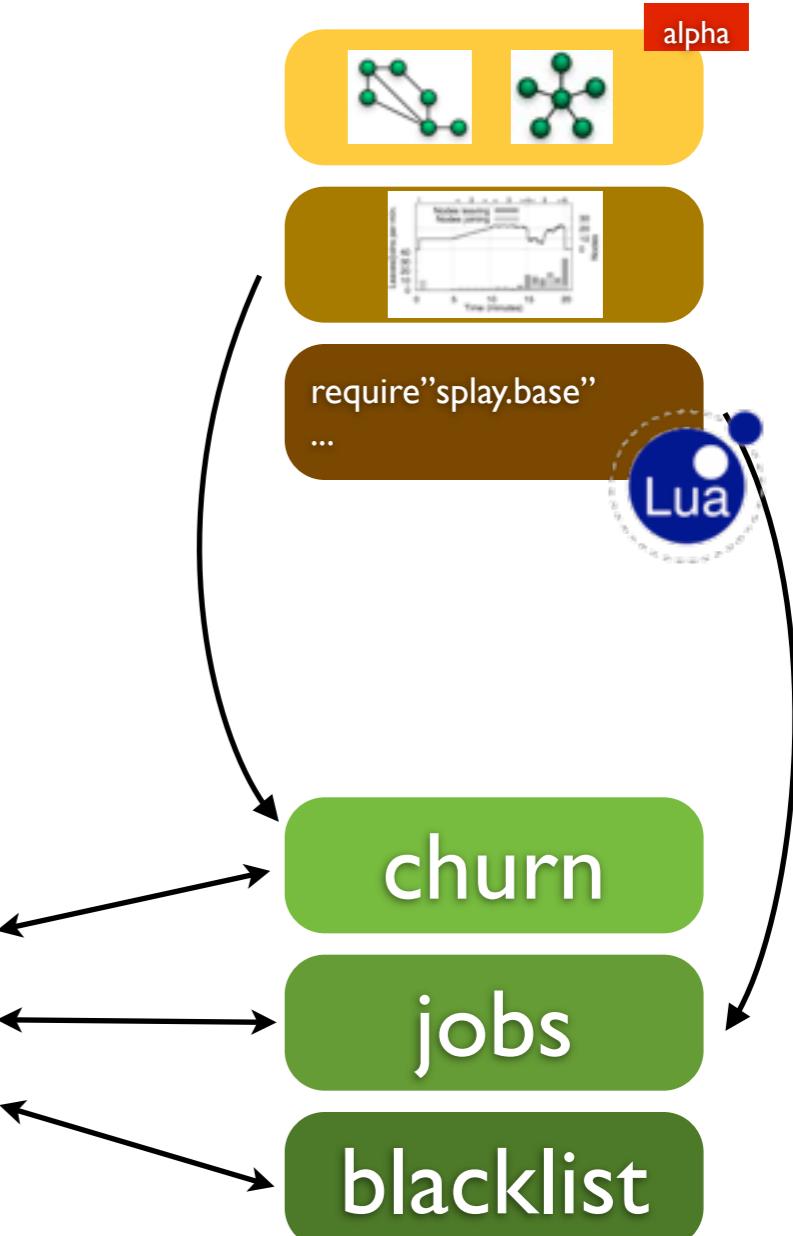
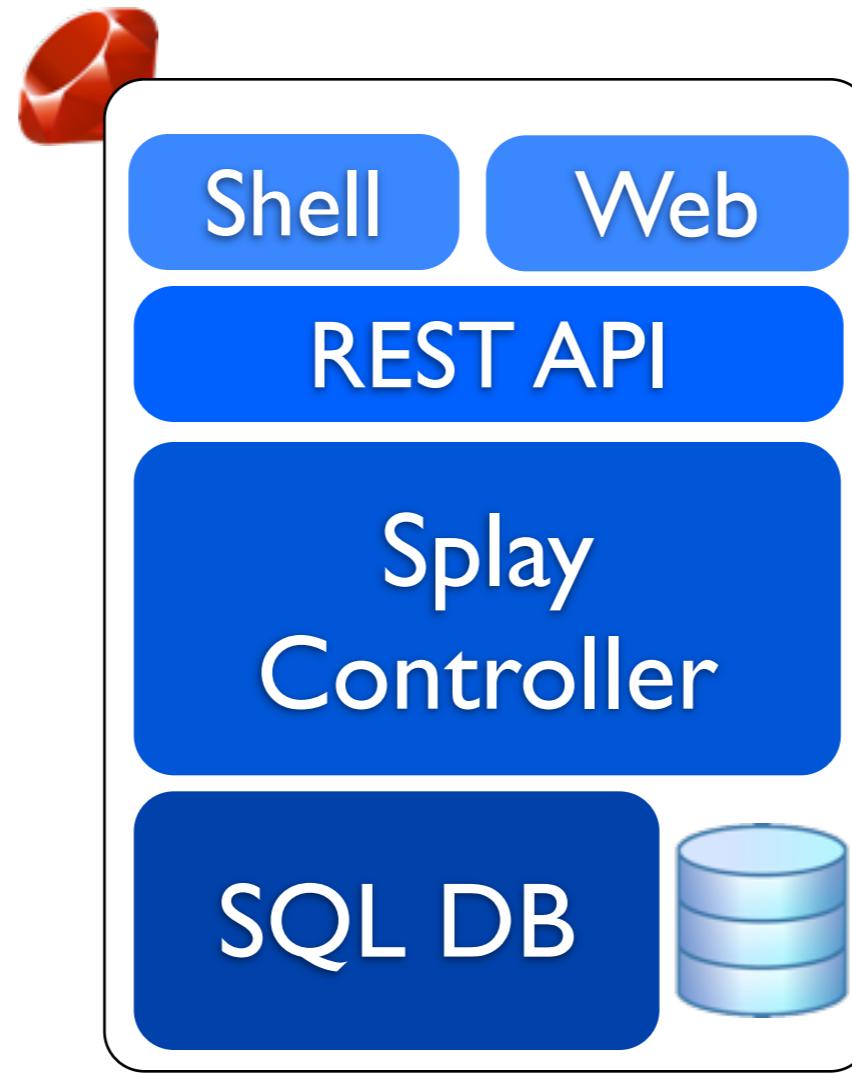
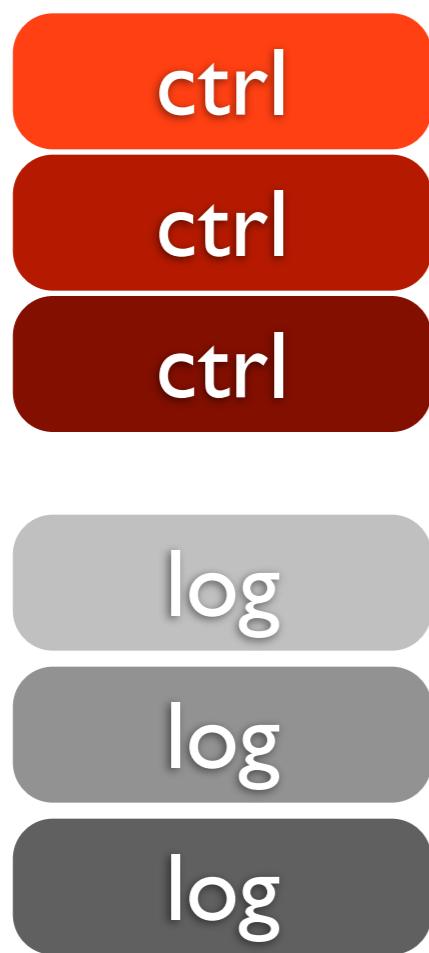
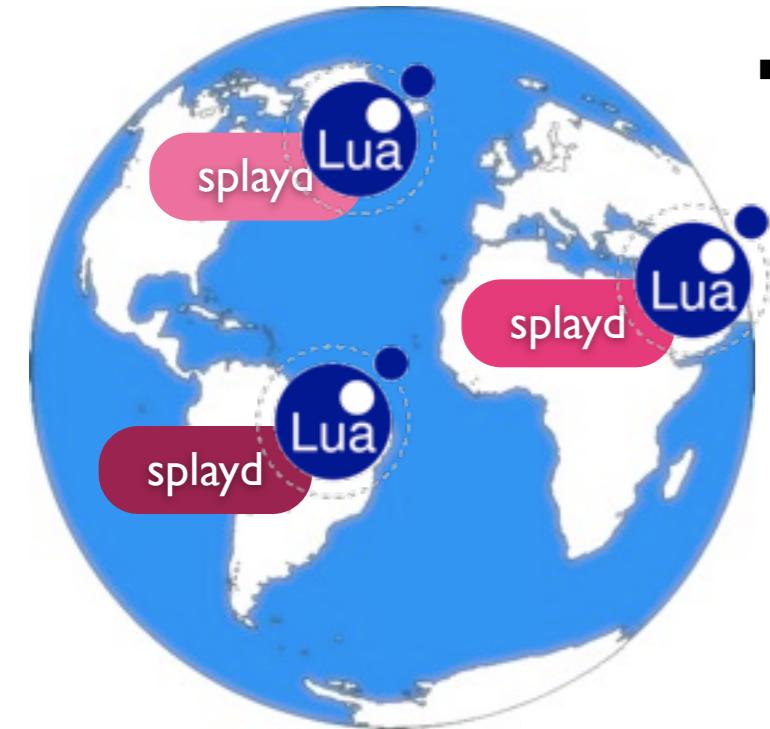
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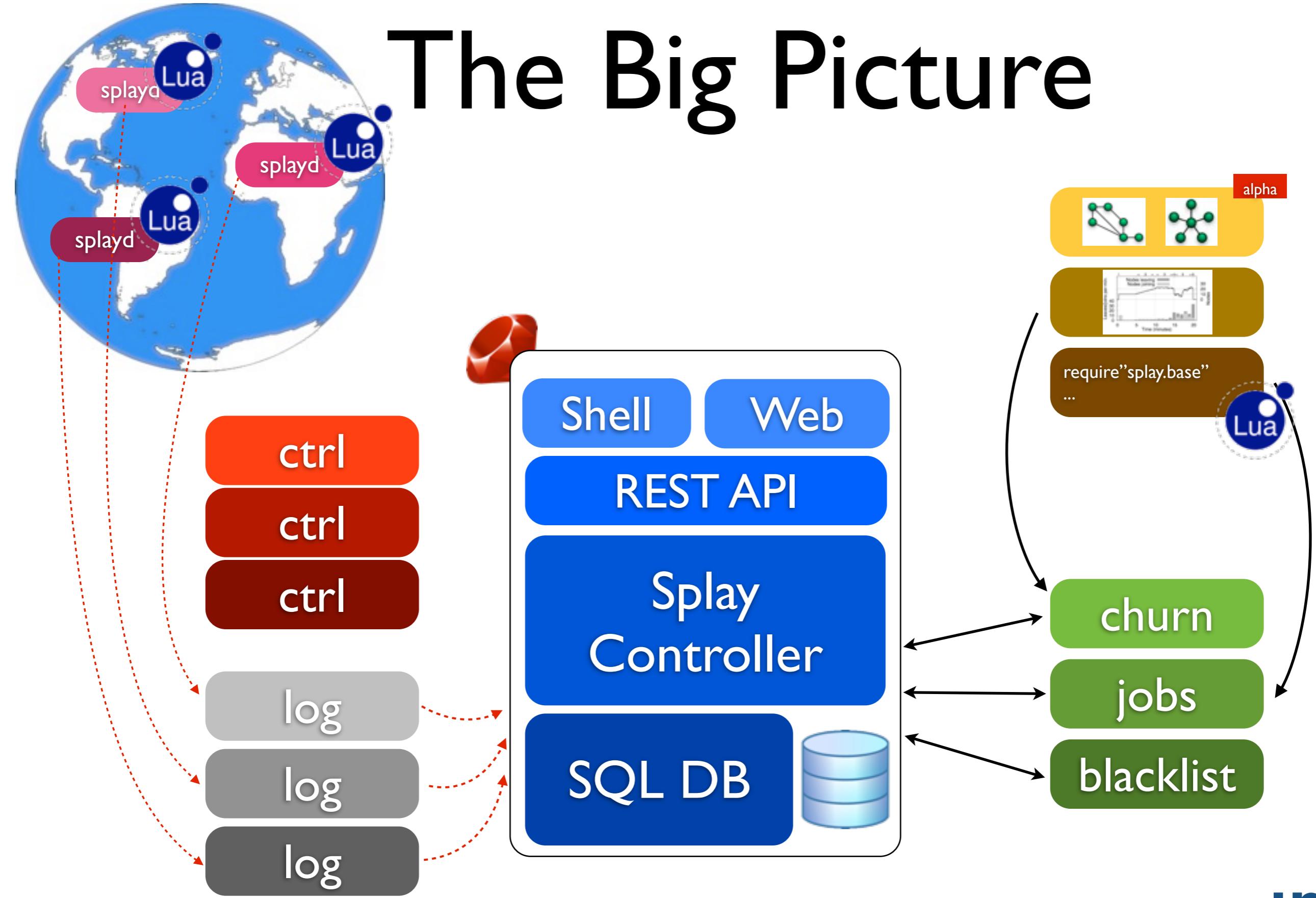
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